## Update on Cheddar : reviewing Multi-Core and ARINC653 scheduling features, software design exploration

P. Dissaux\*, J. Legrand\*, A. Schach\*, S. Rubini+, J. Boukhobza+, L. Lemarchand+, J.P. Diguet+, N. Tran+, M. Dridi+, R. Bouaziz\$, F. Singhoff (speaker)+

\* Ellidiss Technologies
 + Lab-STICC UMR CNRS 6285/UBO
 \$ ReDCAD Laboratory, University of Sfax









# Multiprocessor sheduling analysis with AADLInspector/Cheddar

#### □ SMART project (completed in 2014):

- Define typical multiprocessor architectures AADLInspector should support (pattern)
- □ How to model multiprocessor architectures with AADL
- Choose or design new scheduling analysis methods for those patterns
- □ Prototyping in Cheddar, to be available in AADLInspector

#### □ Main outcomes:

- 1. Implementation of partitioned and global scheduling methods
- 2. Support of shared resources between processing units
- 3. Design of partitioning algorithms

# Typical multiprocessor scheduling analysis: partitioned vs global



- Partitioned scheduling : first assign off-line each task on a processing unit ; each processing unit schedules its own task set.
  No migration. Both on-line and off-line.
- Global scheduling: choose the next task to run on any available processing unit (or preempt if all busy).
  - □ With migration. Fully on-line.

### Typical multiprocessor scheduling analysis: partitioned vs global

#### □ AADLInspector 1.6 :

- Partitioned scheduling only
- Classical policies (fixed priority, EDF, including ARINC 653, ...)
- □ Ravenscar data, data port
- Scheduling simulation & Response time analysis
- Partitioning policies: Best fit, First Fit, Next Fit, GT, SF



#### □ Cheddar 3.1 only (not in Al yet):

- Global scheduling : any uniprocessor policies + specific policies such as EDZL, LLREF, Pfair,
- □ Partitioning policies based on PAES (Pareto Archived Evolution Strategy)
- □ Hardware shared resources support

## Shared resources between processing units



□ Shared resources: Cache units, bus, NoC, ...

- Interferences due to processing units shared resources, make thread WCET (Worst Case Execution Time) difficult to compute
- □ Specific scheduling methods

### **Cache and CRPD**

- □ In fixed priority preemptive scheduling context, tasks can preempt and evict data of other tasks in the cache.
- □ Cache related preemption delay (CRPD): additional time to refill the cache with the cache blocks evicted by the preemption.

#### Some issues:

- □ CRPD is high, non-negligible preemption cost. It can present up to 44% of the WCET of a task (Pellizzoni et al., 2007)
- CRPD is difficult to accurately compute off-line (worst case bound, number of preemption)
- □ Classical scheduling analysis results cannot be applied with CRPD
  - Applying Rate Monotonic priority assignment algorithm may lead to unschedulable task set

□Need new priority assignments taking CRPD into account

## Cache/CRPD-Aware Priority Assignment Algorithms

- Extend Audsley's priority assignment algorithm (Audsley, 1995) to take into account CRPD.
- CRPD-aware priority assignment algorithms (CPA) that assign priority to tasks and verify theirs schedulability.
- 4 algorithms with different levels of schedulability efficiency and complexity.
- □ Implemented into Cheddar 3.1, not available with AADLInspector 1.6

### **Cache-Aware Scheduling Simulation**



#### **Problem Statement:**

- □ Theoretical issues with CRPD : feasibility interval, sustainability
- Various parameters need to be taken into account in scheduling analysis of systems with cache: cache profile, memory layout, CFG

#### ❑ Outcomes:

- □ We have designed a new CRPD computation model, sustainable for L1 instruction cache. Feasibility interval proved.
- □ Extending Cheddar to model cache/cache access profile

## **Summary**

- 1. Multiprocessor scheduling analysis features
- 2. Software design space exploration : partitioning with competing objective functions

## Cheddar & partitionning with competing objective functions

#### □ Problem statement :

□ Performances (scheduling), is not the unique concern

- Trade-offs with several competing criteria/objective functions such as performances vs safety vs security
- □ How to do partitionning in this context ?
- □ PAES helps ? PAES with Cheddar ?

#### □ Small example to illustrate, assume:

- □ A system running several sub-programs (i.e. functionnal units)
- □ Subprograms may shared resources (compliant with Ravenscar)
- □ How to assign subprograms to threads

## From the functional specification to a software architecture



# Competing objective functions in software design space exploration



Explore several assignment solutions

Select assignment solutions that meet at **best the trade**offs between number of preemptions and laxities

## **PAES : a multi-objectives metaheuristic**

### □ Basic steps of PAES algorithms:

Mutate a solution to generate a new candidate: small change to move from a solution to a nearby neighbour

**Evaluate** the mutated solution (conflicting objective functions)

Update non-dominated solutions set (i.e. archive)

Select new solution for next iteration : mutated or current solution

#### Pareto Front: final set of non-

dominated solutions

1

2

3

4

 Solutions A dominates solution C because it is better than C for all objectives



#### **PAES-based partitioning**



## Competing Performance Criteria in the Software Design Space exploration

## □ Examples of investigated trade-offs with competing objectives functions such as:

□Min (#premptions)

□Max (laxities)

□Min (Ravenscar data blocking time)

**D**...

⇒Performance competing objectives functions only

## □ How to be sure that objective functions are competing?

### Conclusion

#### Multiprocessor scheduling analysis of AADLInspector & Cheddar:

Bunch of classical partitioned vs global scheduling algorithms

□Shared hardware resources: cache, NoC

#### □ Multi-objective partitioning

PAES based, for Ravenscar compliant architecture
 Safety & performance & security objective functions
 Follow Security annex